

Spreading Metrics and Optimal Linear Arrangement

OLA: Given a graph $G(V, E)$ with edge weights $w_e \geq 0$

Find a bijection $f: V \rightarrow \{1, 2, \dots, n\}$ such that

$$\sum_{e=(u,v)} w_e |f(u) - f(v)| \text{ is minimized}$$

map a graph vertices to a straight line
without distorting the edge lengths too much.

$$d_{uv} = |f(u) - f(v)|$$



Spreading
Metric

Constraint
└→

$$\min \sum_{e=(u,v)} w_e d_{uv} \text{ s.t.}$$

$$\forall S \subseteq V \text{ and } u \notin S, \sum_{v \in S} d_{uv} \geq \frac{1}{4} |S|^2$$

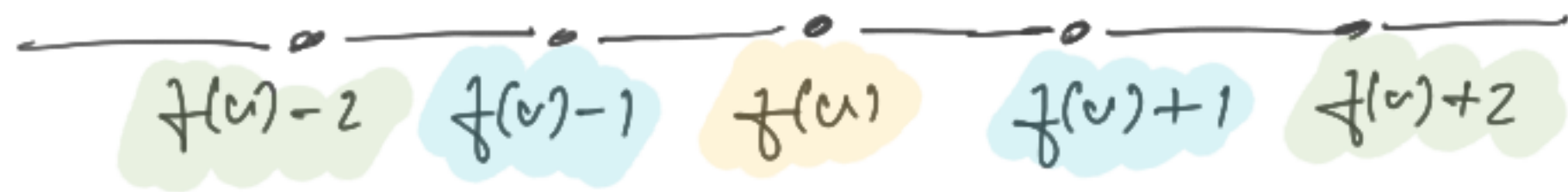
makes
 d_{uv} a
metric

$$\left\{ \begin{array}{ll} d_{uv} = d_{vu} & \forall u, v \in V \\ d_{uv} \leq d_{uw} + d_{wv} & \forall u, v, w \in V \\ d_{uv} \geq 1 & \forall u \neq v \\ d_{uu} = 0 & \forall u \end{array} \right.$$

└→ Above constraints make it a Spreading Metric

Claim: Above LP is a relaxation of OLA

- Given a mapping $f: V \rightarrow \{1 \dots n\}$, $d_{uv} = |f(u) - f(v)|$
- Consider $S \subseteq V$ and $u \notin S$
 - at most 2 vertices in S at dist 1 from u
 - at most 2 vertices in S at dist 2 from u
 - $\vdots$
 - at most 2 vertices at dist $|S|/2$ from u



$$\therefore \sum_{v \in S} d_{uv} \geq 2(1 + 2 + \dots + |S|/2) \geq \frac{|S|^2}{4}$$

Hence a solⁿ to OLA is feasible for this LP

- Solving this LP: use Ellipsoid method with a separation oracle for the Spreading-Metric constraints.

- Given a candidate solⁿ $\{d_{uv}\}$, need to check if $\exists S \subseteq V$ and $u \notin S$ st $\sum_{v \in S} d_{uv} < |S|^2/4$

- For each u , sort $V - \{u\}$ by distance from u
 $v_1 \ v_2 \ \dots \ v_{n-1}$

- Claim: If $\exists S \ni u$ st $\sum_{v \in S} d_{uv} < |S|^2/4$
then $\sum_{i=1}^{|S|} d_{uv_i} < |S|^2/4$

Because $\{v_1, \dots, v_{|S|}\}$ and S have same size
and distances to u can only be greater

- Hence enough to test $\{v_1\}, \{v_1, v_2\}, \dots, \{v_1, v_2, \dots, v_{n-1}\}$

Obs: If d is a Spreading Metric, for any $S \subseteq V$ and $u \notin S$

$\exists v \in S$ st $d_{uv} \geq |S|/4$

- $\sum_{v \in S} d_{uv} \geq |S|^2/4 \Rightarrow \text{avg dist to } u \geq |S|/4$

Approx Algo:

- Solve LP to get a Spreading Metric (V, d)
- Construct a Tree Metric (V', T) where
 - $V \subseteq V'$ is at the leaves of T
 - $\sum_{u, v \in V} w_{uv} T_{uv} \leq O(\log n) \sum_{u, v \in V} w_{uv} d_{uv}$
 - If $z \in V$ is in the smallest subtree of T containing $u, v \in V$, then $T_{uv} \geq T_{uz}$
- Map V , the leaves of T , to $\{1, 2, \dots, n\}$ from Left to Right

↪ given f



Th^m: This is a $O(\log n)$ approx to OLA

- Pf:
- Let $f: V \rightarrow \{1 \dots n\}$ be the computed mapping
 - For $e = (u, v)$ consider the smallest subtree of T containing u and v
 - leaves of this subtree are mapped to consecutive integers in a range $[a, b] \Rightarrow b - a + 1$ leaves
 - $|f(u) - f(v)| \leq b - a$
 - $S = \text{leaves except } u, |S| = b - a$
 - By Obs, $\exists z \in S$ st $d_{uz} \geq \frac{b-a}{4}$
 - By Tree Metric properties, $T_{uv} \geq T_{uz} \geq d_{uz} \geq \frac{b-a}{4}$
 - $\therefore |f(u) - f(v)| \leq 4 T_{uv}$

$$\begin{aligned}
 \therefore \sum_{e=(u,v)} w_e |f(u) - f(v)| &\leq \sum 4 w_e T_{uv} && \text{(from above)} \\
 &\leq O(\log n) \sum w_e d_{uv} && \text{(by Tree Metric property)} \\
 &\leq O(\log n) \text{OPT}_{\text{OLA}} && \text{(by LP relaxation)}
 \end{aligned}$$

Cost of the
approx solⁿ



Thm: Given metric (V, d) and costs $c_{uv} \forall u, v \in V$,
in deterministic polynomial time we can compute

a tree metric (V', T) st

- $V \subseteq V'$ are leaves of T

- $T_{uv} \geq d_{uv} \forall u, v \in V$

- $\sum_{u,v} c_{uv} T_{uv} \leq O(\log n) \sum_{u,v} c_{uv} d_{uv}$

- T is rooted and for any $z \in V$ that
belongs to the smallest subtree of T containing $u, v \in V$
we have $T_{uv} \geq T_{uz}$

$$\forall u, v \quad \mathbb{E}[T_{uv}] \leq O(\log r) d_{uv}$$

